

The P vs. NP problem

**Efficient computation, Internet security,
and the limits of human knowledge**

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Clay Math Institute Millennium Problems - \$1M each

- Birch and Swinnerton-Dyer Conjecture
- Hodge Conjecture
- Navier-Stokes Equations
- P vs. NP



~~Poincaré Conjecture~~

- Riemann Hypothesis
- Yang-Mills Theory

Scientific / Mathematical/ Intellectual / Computational problems



NP: Problems we want to solve/understand

The diagram consists of two white ellipses on a starry background. The larger outer ellipse is labeled 'NP: Problems we want to solve/understand'. Inside it, towards the bottom right, is a smaller ellipse labeled 'P: Problems we can solve/understand'. This visualizes that P is a subset of NP.

P:
Problems we can solve/understand

P=NP? - limits on human knowledge

PLAN

- Computation is everywhere
- Algorithms: language of computation
- Efficient algorithms: P
- Efficient verification: NP
- NP-completeness
- Implications

Computation

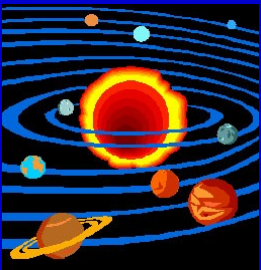
Mathematics

$$X^n + Y^n = Z^n$$

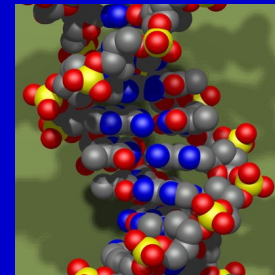
Computer



Computation



Physics



Biology

everywhere

Computation: every process
which is a
sequence of *simple, local* steps,
that
we want to perform, or
understand

Variety of natural phenomena
and intellectual challenges, each
with an essential computational

1 month



input



2 pm

**Fetal
development**

**Weather
evolution**

3 months



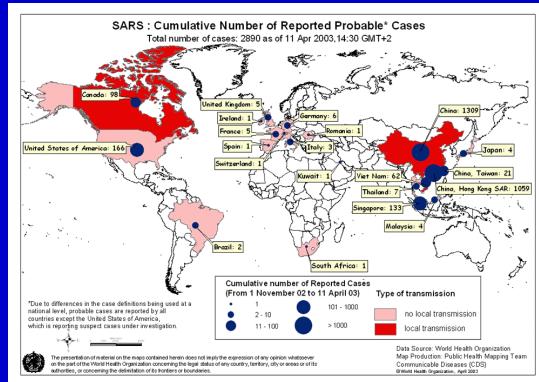
output



4 pm

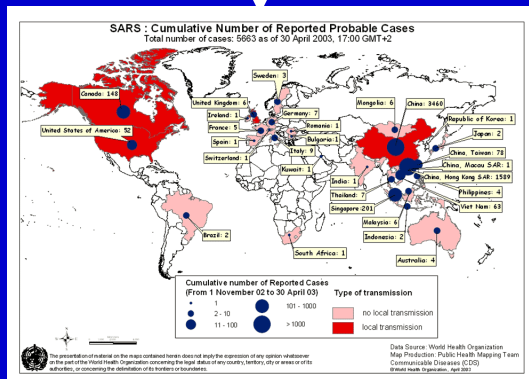
**Nature computes !
Can we simulate/predict?**

4/11/03

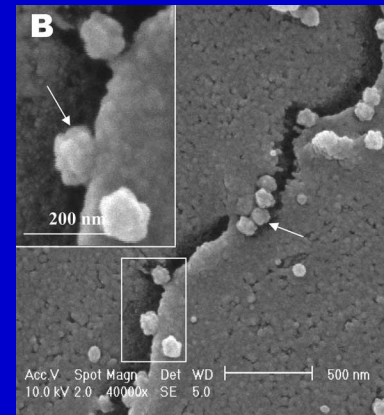


**SARS infection
(in the world)**

4/30/03

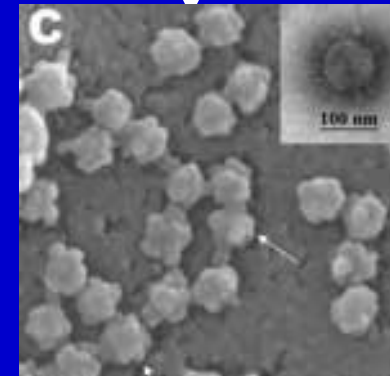


Will the epidemic spread, or die out?



+15h

**SARS infection
(in the cell)**



24h

$$X^2 + Y^2 = Z^2$$

Solving
equations

$$X=3 \ Y=4 \ Z=5$$

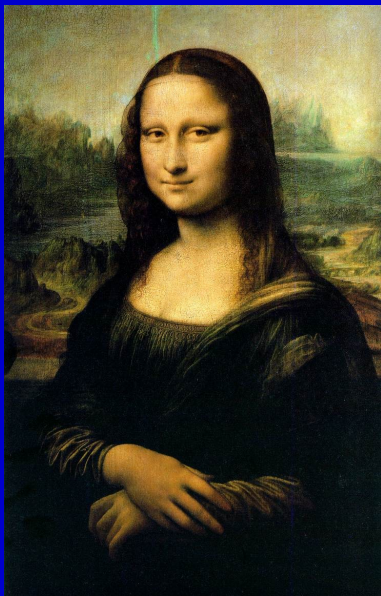
$$X^n + Y^n = Z^n \ n>2$$

Proving
theorems

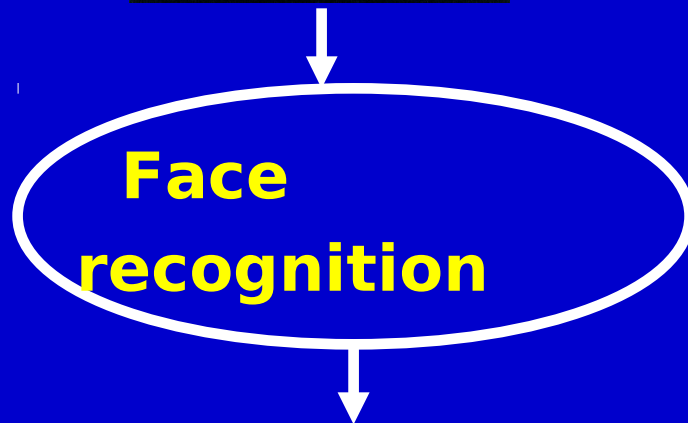
Theorem: no solution!
Proof does not fit on
this slide (200 pages)



Computations in Mathematics

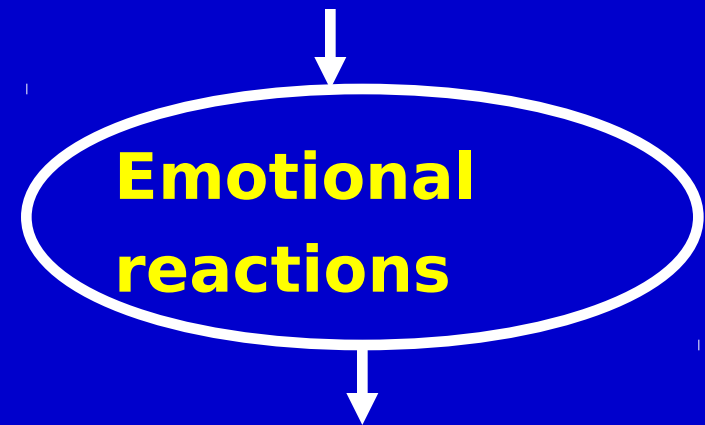


Nearly 10,000 reported killed by China quake
Rain hampering rescue efforts in worst-hit area
Nearly 900 children buried when a school building collapses, 50 bodies found
7.9 magnitude quake is felt throughout much of China



**Face
recognition**

"Mona Lisa"



**Emotional
reactions**

Sadness

The subconscious brain computes

Beauty from computation



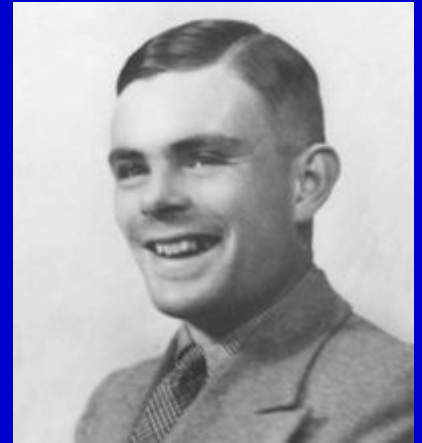
Seashells compute

**How to describe
computation?**

**The language of
Algorithms**

Father of Computing

Alan Turing 1912-1954



1936: “On computable numbers, with an application to the entscheidungsproblem”

- Formal definition of **algorithm** (Turing machine)
- Seed of the computer revolution
- **Church-Turing Thesis:** everything that nature computes, can be emulated on a Turing machine

ALGORITHM (informal)

Step-by-step, **local**,
simple, mechanical
procedure.

Halts in **finite** time
for every input.

		1	1	1			
		1	2	3	4	5	
			6	7	8	9	
		1	9	1	3	4	

Example: Addition algorithm (informal)

1. Scan column. If empty, stop.
2. Add digits. Write answer, remember carry.
3. Move one column left, write carry.
4. Go to 1

~~Finite description vs. Infinite # inputs~~

Limits of Knowledge I

Unsolvable

Turing (& Godel):
Given a computer
program, does it
always halt?

Mattiasevich: Given
an
equation, does it
have an integer
solution?

Conway: Given a

Solvable

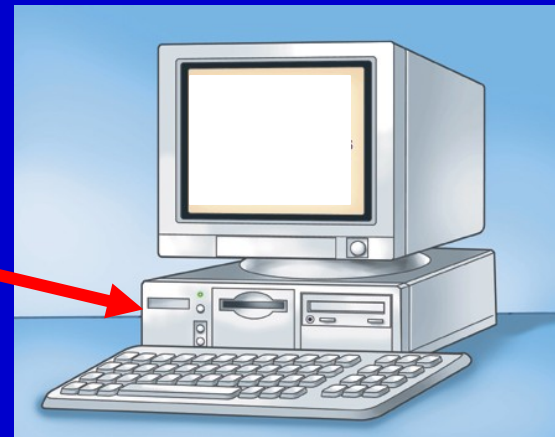
When?



Computational
Complexity
Theory

Efficiency of an algorithm -
asymptotic analysis:
Number of basic steps,
for larger and larger
inputs.

input



Rubik's cube

How many steps to
solve..



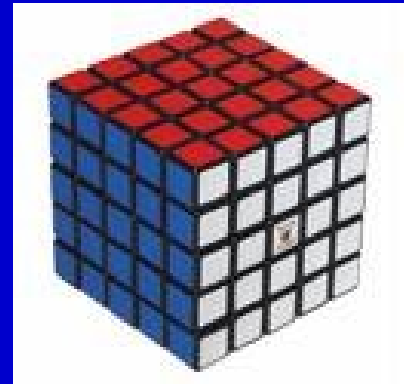
2



3



4



5

...

Sudoku

How long does it take you to solve...

		8	6					
							6	
			4	8			2	3
		5		9				8
	4	9				2	1	
2				4		7		
3	6			2	9			
	1							
					5	1		

3

1			2	3	4			12		6				7	
		8				7			3			9	10	6	11
	12			10			1		13		11			14	
3			15	2			14				9			12	
13				8			10		12	2		1	15		
	11	7	6				16				15			5	13
			10		5	15			4		8			11	
16			5	9	12			1						8	
	2						13			12	5	8			3
	13			15		3			14	8		16			
5	8			1				2				13	9	15	
		12	4		6	16		13			7				5
	3			12				6			4	11			16
	7			16		5		14			1			2	
11	1	15	9			13			2				14		
	14				11		2			13	3	5			12

4

Sudoku

	y				b		a	c	x	n			h			t		f	l			d	e		
	t		s	u	j	h	v				d	q	c			o	k		b	n		a	w	p	
w	h	e	m	a	n		l	u	k	p		r	y			s	x	d	q	c	o	j		i	b
b		j		p	s			t				i	m			v	n	g	h	a	q		r	x	y
x	o	l	d		i	p			r	e			f			u	j	w	y	m		h		s	c
	w	q	u	j			i		e		x	b	o	m		a			n	h	k	c	s		
n	c				w	x	u	s	f		q		l			e	m	k	v				j	a	
a	i		x	f	c	l			m		v	k	w		q			j		d	g		b	h	
s			v		h	k	p	o	b	u	f	j	n				t		d	i	m		r	q	
	b	d		m	r	v			j		h	p				o	g	y	w			t		u	
y	p		e	l	a	m		v	h	o	b		x	i		t	s	q	u	w	g	r	c	d	k
	q	g	j		e		s	r		h	c					f	k			x		y	l	a	o
		u	t	k		n	o		l		r	m	q	y		b	a	v	j		i	p	h		
	x	r		w	p		y	k	i		l	e	j					m			t	q	v		u
	s		n	b	q	c		g	w	k	a	u	t	p	y		o		r	x		j	m		
j	n	s	q	v	x	y	h		u	t	p	o	g	l	m		f	d			w	i	k	r	
u		w	b	t	l	e	r	p	o	m		c	d	f	k	v					s	q			
d			h		m	s	c	f		q	j		k	n	g	w		b			l	v	u		e
			o	e	d	i	k	n	q		w		u		j	a	l				h		b	p	m
l	k				v	j	t	w		a	s	h					u	r	q	c	d	f		n	
			g	d	y	r	w			c		l	i		n	p	v	a	f	e			q		
	v	x	p	o		t	b			d	n	f			w			g		s	a	h	y	i	
i		k	w	c	g	q	x	h					a	u	l	d	e		s			m	f	v	
		a	y	r		d	f	e	n	x	k		s	h			b		u		p				
q	l		f	s			m	i	v			w			h		x	t	y				c	d	

Efficiency of the addition algorithm

5 DIGITS
30 STEPS

1. Add digits. Write answer, retain carry.
2. Move one column left, write carry.
3. Scan column. If empty, stop.
4. Go to 1

12345
+6789

10 DIGITS

60 STEPS

123456789

+987654321

6 basic steps per column

20 DIGITS

72635273545786043726

120 STEPS

+53827484732625435473

50 DIGITS
300 STEPS

47563739203487456438992305757328576452364568456465744576

Is there a faster algorithm?

No!

2

N DIGITS

6N STEPS

Solving is as fast as reading the input

Efficiency of the multiplication algorithm

5 DIGITS

25 STEPS

10 D

100

Grade-school multiply algorithm

Diagram illustrating the construction of a 2D array from a 1D array. The top row shows a 1D array of 15 elements, with the first element 'X' and the rest '*'. A horizontal line separates this from the 2D array below. The 2D array is a triangular matrix of '*' characters, with the label n^2 to its left. The bottom row of the 2D array is a 1D array of 15 elements, all '*'. A horizontal line separates the 2D array from this bottom row.

12345

x6789

123456789

X
987654321

400 STEPS

72635273545786043726

x53827484732625435473

50 DIGITS

2500 STEPS

47563739203487456438992305757328576452364568456465744576

98656092843467546234868431987543210979832865874134653472

I there a faster algorithm? **Yes!**

N DIGIT

sec But

But not as fast as addition

in a

Efficiency of a factoring algorithm

?

×

?

= 147,573,952,588,676,412,927

Find nontrivial factors of a number A

N DIGITS

10^{N^2} STEPS

Brute force factoring algorithm

Input: A

- For $B = 2, 3, \dots, \sqrt{A}$ do:
- If B divides A, return B, A/B

Very slow! 1000 digits → sun will die before finishing

Is there a faster algorithm?

Yes, but still extremely slow!

Which problems are hard to solve?

Addition & Multiplication: Easy

Is Factoring hard ?

**Finding efficient algorithms, or
proving that no such algorithms
exist:**

Bread and butter of our field

Cobham, Edmonds
Rabin ~1965

The class P

**All problems having an efficient
(polynomial time, e.g. n , n^2)
algorithm**

like Addition and Multiplication

**Many practical interesting problems
in P**

Efficient algorithms -

Drivers of invention & industry

Who were

**Edison ? Marconi ? Guttenberg ?
Stevenson ?**

**Light bulb Radio Printing press Steam
engine**

Shortest path

1959
Google Maps

MAPQUEST.

Network flows
Internet routing
Dynamic
Programming



```
define Dijkstra(Graph G, Node s)
  S := {}
  Q := Nodes(G)
  while not empty(Q)
    u := extractMin( Q )
    S := S ∪ u
    for each node v in neighbors( u )
      if d(u) + w(u,v) < d(v) then
        d(v) := d(u) + w(u,v)
        pi(v) := u
```

Distance (Delhi, Bangalore)
Path (Delhi, Bangalore)

Pattern matching

Knuth-Morris-Pratt
Boyer-Moore 1977

Spell checking

Text processing

Genome

Molecular Biology



Text CAUCGCGCUUCGC
Pattern CGC



```
algorithm kmp_search:
  input: T (text), P (pattern sought)
  define variables:
    m ← 0, i ← 0, M (the table)
  while m + i is less than length of T, do:
    if P[i] = T[m + i], let i ← i + 1
    if i = length of P then return m
    otherwise, let m ← m + i - M[i],
    if i > 0 let i ← M[i]
```



Text CAUCGCGCUUCGC
Location X X X

Fast Fourier Transform (FFT)

Cooley-Tukey 1965

Gauss 1805

Audio processing
Image processing
Tomography, MRI
Fast multiplication
Quantum algorithms

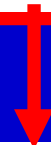


$T(0), T(1), T(2), \dots, T(N)$



RECURSIVE-FFT(a)

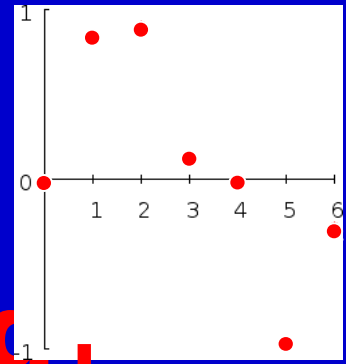
```
1   $n \leftarrow \text{length}[a]$ 
2  if  $n = 1$ 
3    then return  $a$ 
4   $\omega_n \leftarrow e^{2\pi i/n}$ 
5   $\omega \leftarrow 1$ 
6   $a^{[0]} \leftarrow (a_0, a_2, \dots, a_{n-2})$ 
7   $a^{[1]} \leftarrow (a_1, a_3, \dots, a_{n-1})$ 
8   $y^{[0]} \leftarrow \text{RECURSIVE-FFT}(a^{[0]})$ 
9   $y^{[1]} \leftarrow \text{RECURSIVE-FFT}(a^{[1]})$ 
10 for  $k \leftarrow 0$  to  $n/2 - 1$ 
11   do  $y_k \leftarrow y_k^{[0]} + \omega y_k^{[1]}$ 
12      $y_{k+(n/2)} \leftarrow y_k^{[0]} - \omega y_k^{[1]}$ 
13      $\omega \leftarrow \omega \omega_n$ 
14 return  $y$ 
```



$$T_N(x) = \sum_{n=0}^N a_n \cos(nx) + i \sum_{n=0}^N a_n \sin(nx)$$

Error correction

Reed-Solomon decoding



Petersen 60

Berlekamp-Massey 6

CDs



DVDs



Satellite communication

Cell phone communication



INPUT: a binary sequence $S = S_0, S_1, S_2, \dots, S_n$.

OUTPUT: the complexity $L(S)$ of S , $0 < L(S) < N$.

1. Initialization: $C(D) := 1$, $L := 0$, $m := -1$, $B(D) := 1$, $N := 0$.

2. While $(N < n)$ do the following:

2.1 Compute the next discrepancy d .

$$d := (S_N + \sum c_i S_{N-i}) \bmod 2.$$

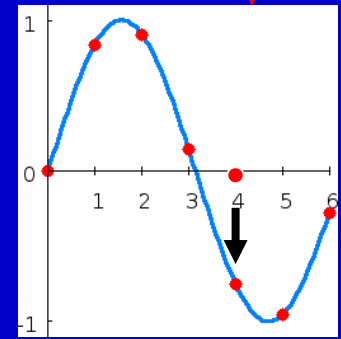
2.2 If $d = 1$ then do the following:

$$T(D) := C(D), C(D) := C(D) + B(D) \cdot D^{N-m}.$$

$$\text{If } L < N/2 \text{ then } L := N+1-L, m := N, B(B) := T(D).$$

2.3 $N := N+1$.

3. Return(L).



**Unsolvab
le**

**Solvabl
e**

**Shortest
Path**

**Pattern
Matching**

P

FFT

**Error
Correction**

Multiplication

Addition

Cobham, Edmonds

Rabin ~1965

The class **P**

**All problems having an efficient
(polynomial time) algorithm**

Many interesting problems in **P**

Are **all interesting problems in **P**?**

What are “interesting” problems?

Search problems

Short Path: FIND short path from Princeton to LA

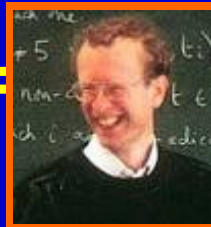
Pattern Matching: FIND CGC in
CAUCCGCCUCCG

Easy What is common to all these problems?
Hard In all, solutions are easy to **check & verify!**

Factoring: FIND factors of 147,573,952,588,676,412,927
= 193,707,721 ×

761,838,257,287

Theorem Proving: FIND



Lemma...Proof...Lemma..Proof..

200-page proof of the
man

		8	6				
						6	
			4	8		2	3
		5		9			8
	4	9				2	1
2				4		7	
3	6			2	9		
	1						
				5	1		

9	2	8	6	1	3	4	5	7
4	7	3	9	5	2	8	6	1
1	5	6	4	8	7	9	2	3
7	3	5	2	9	1	6	4	8
6	4	9	7	3	8	2	1	5
2	8	1	5	4	6	7	3	9
3	6	7	1	2	9	5	8	4
5	1	2	8	7	4	3	9	6
8	9	4	3	6	5	1	7	2

Sudoku: FIND solution of

The class NP- problems like **FIND**: needle in a haystack



May be **hard** to *find* Always **easy** to *verify*

Cook & Levin 1971

Gödel 1956

The class NP

All problems having efficient
verification

algorithms of given solutions

For every such problem, **finding** a
solution (of length **n**) takes $\leq 2^n$
steps: try all possible solutions &
verify each.

Can we do better than “brute
force” ?

**Unsolvab
le**

**Solvabl
e**

**Integer
Factoring**

**Shortest
Path**

**Pattern
Matching**

**Solving
Sudoku**

**Theorem
Proving**

P

FFT

**Error
Correction**

Multiplication

Addition

N

P

P versus NP

P: Problems for which solutions can
be efficiently *found*

NP: Problems for which solutions can
be efficiently *verified*

Conjecture: $P \neq NP$

[finding is much harder than verification]

“**P=NP?**” is a central question of
math, science & technology !!!

What is in NP?

Mathematician: Given a statement, *find* a proof

Scientist: Given data on some phenomena,
find a theory explaining it.

Engineer: Given constraints
(size, weight, energy)
find a design (bridge, medicine,
phone)

**In many intellectual challenges, *verifying*
that we found a good solution is an easy
task !**

Universality: NP-completeness

Are SuDoku, Theorem Proving, Factoring hard?

These problems are intimately related!!

Theorem: If SuDoku is easy then

- Theorem proving is easy
- Factoring is easy

Proof: SuDoku is NP-complete

$P = NP$ iff SuDoku has an efficient algorithm

SuDoku solver can solve any NP problem

Universality: NP-completeness

NP-complete problems:

If one is easy, then all are!

If one is hard, then all are!

SuDoku: NP-complete

Thm proving: NP-complete

Integer factoring: we don't know

**Unsolvab
le**

**Solvabl
e**

**NP-
complete**

Solving Theorem Proving
Sudoku

Integer
Factoring

Shortest
Path

Pattern
Matching

P

FFT

Error
Correction

Multiplication

Addition

**N
P**

Universality: NP-completeness

NP-complete problems:

If one is easy, then all are!

If one is hard, then all are!

SuDoku: NP-complete

Thm proving: NP-complete

Integer factoring: we don't know

**Thousands of NP-complete problems
known in Math, Biology, Physics,
Economics,....**

Protein Engineering vol. 7 no. 9 pp. 1059-1068, 1994

The protein threading problem with sequence amino acid interaction preferences is NP-complete

Richard H. Lathrop

Economic Theory vol. 23, no. 2 , pp. 445-454, 2004

Finding a Nash equilibrium in spatial games is NP-complete

R. Baron, J. Durieu, H. Haller and P. Solal

[math.GR] arXiv:0802.3839v1

Quadratic equations over free groups are NP-complete

O. Kharlampovich, I.G. Lysenok, A G Myasnikov,
N. Touikan

**NP-completeness: sign of structural
“nastiness”.**

Potential guide to better models and

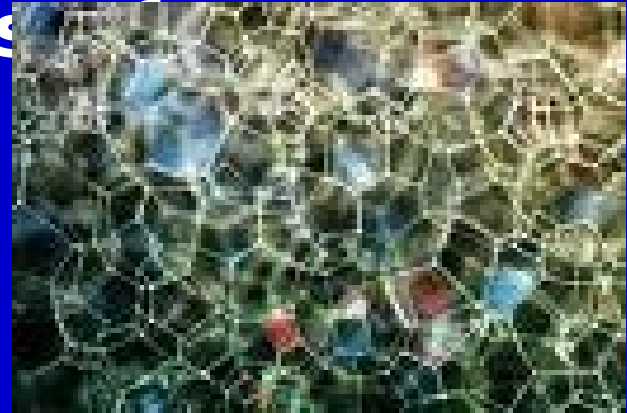
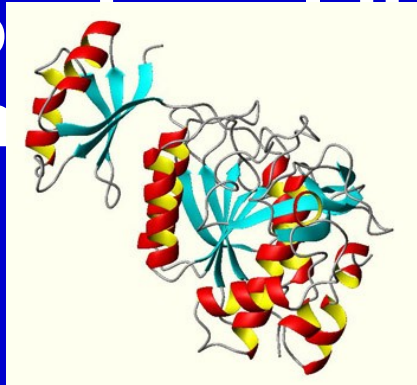
$P \neq NP$ as a law of nature

The following problems are NP-complete

Biology: Minimum energy **Physics:**

Minimum

Protein Folding
Foam



Economics: Nash Equilibrium in strategic games

What is efficient computation?

Church-Turing Thesis:

efficiently?

Every *reasonable* process, can be

- Adding random bits
simulated by a Turing machine

Theorem [Blum-Micali, Yao, Nisan-Wigderson, Impagliazzo-Wigderson]

If “ $P \neq NP$ ”, randomness add no power!

- Adding quantum bits

Theorem [Shor]

Positive consequences of $P \neq NP$

$P \neq NP$ Some of the problems we want to solve are hard. Are hard problems useful?

Cryptography: If **Factoring** is hard then:

- Encryption
- Electronic commerce
- Digital signatures
- On-line shopping

Things we didn't cover

- How to prove NP-completeness
- Attempts to prove $P \neq NP$ and restricted lower bounds
- Other resources (space, parallelism communication) and complexity classes
- Other modes of computation (average-case, approximate,...)
-

